

CS1800  
Discrete Structures  
Fall 2017

Lecture 24  
11/1/17

## Last time

- Expectation

## Today

- Applications of Expectation
  - Variance & Standard Deviation
  - Entropy

## Next time

- Conditional Prob.
- Bayes Theorem

# Variance

Consider heights

## case 1

$$4'10'' \ 5' \ 5'2''$$

$$E[x] = 5'$$

## case 2

$$4' \ 5' \ 6'$$

$$E[x] = 5'$$

## case 3

$$3' \ 5' \ 7'$$

$$E[x] = 5'$$

How to measure "variability"

3 ways

$$\begin{aligned} \textcircled{1} \quad E[y] &= \sum_{w \in \Omega} y(w) \cdot p(w) \\ &= (-12") \cdot 1/3 + (0") \cdot 1/3 + (+12") \cdot 1/3 \\ &= 0 \end{aligned}$$

X neg. & pos. cancel

$$\textcircled{2} \quad y = |x - E[x]|$$

- mean absolute deviation

$$\textcircled{3} \quad y = (x - E[x])^2$$

- variance

$$\begin{aligned} \textcircled{2} \quad E[y] &= |-12| \cdot 1/3 + |0| \cdot 1/3 + |+12| \cdot 1/3 \\ &= 12 \cdot 1/3 + 0 \cdot 1/3 + 12 \cdot 1/3 = 8 \text{ inches} \end{aligned}$$

$$\begin{aligned} \textcircled{3} \quad E[y] &= E[(x - E[x])^2] = \text{case 1} \\ &= (-2)^2 \cdot 1/3 + (0)^2 \cdot 1/3 + (2)^2 \cdot 1/3 \\ &= 4/3 + 0 + 4/3 = \boxed{8/3 \text{ inches}^2} \end{aligned}$$

variance  $\sigma^2$

Take square root, get back inches ...

$$\sigma^2 = \text{Var}(X) = E\{(X - E[X])^2\}$$

$$\Rightarrow \sigma = \sqrt{\text{Var}(x)} = \sqrt{E[(X - E[X])^2]} \quad - \text{back in units originally measured.}$$

e.g.

$$\text{case 1 : } \sigma^2 = \text{Var}(X) = 8/3 \text{ inches}^2$$

$$\sigma = \sqrt{\text{Var}(X)} = \sqrt{8/3 \text{ inches}^2} = 1.63 "$$

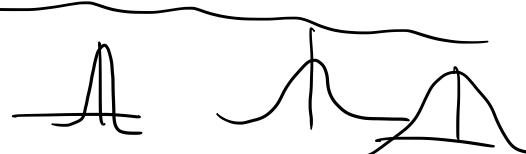
$$\text{case 3 : } \dots \sigma = 19.6 "$$

Exam

$$E[X] = 76$$

$$\sigma = 15 \quad \pm 2/3 \sigma - \text{quartiles}$$

$$2/3 \sigma = 10 \text{ pts}$$



## Entropy

- Consider 8 letters only  $\{A, B, C, D, E, F, G, H\}$

- Need 3-bits

$$A \rightarrow 000$$

$$B \rightarrow 001$$

$$H \rightarrow 111$$

} fixed length code

- Suppose must encode  $n$ -thess, need:  $\lceil \log_2 n \rceil$

- Efficiency of code is measured in bits-per-character on average, BPC.  $BPC = 3$

- Variable length code: assign short codes to frequent letters long codes to infrequent letters.

$$BPC = \sum_{w \in \Sigma} X(w) \cdot p(w) = \sum_i l_i \cdot p_i$$

Why must we have long codes to make up for short codes?

Analog to PHP: Kraft's Inequality

Example:

$$\begin{array}{ll} A \ B \ C \ D \ E \ F \ G \ H & BBB \rightarrow 010101 \\ 00 \ 01 \ 010 \ 011 \ 100 \ 101 \ 110 \ 111 & CF \rightarrow 010101 \end{array}$$

↑  
code length  
for  $i^{th}$  letter      ↑ probability of  $i^{th}$  letter

- Suppose we have following skewed distribution

$$P_i = \left( \frac{A}{4}, \frac{B}{4}, \frac{C}{8}, \frac{D}{8}, \frac{E}{16}, \frac{F}{16}, \frac{G}{16}, \frac{H}{16} \right)$$

(00, 01, 100, 101, 1100, 1101, 1110, 1111)

$ADF \rightarrow 001011101$

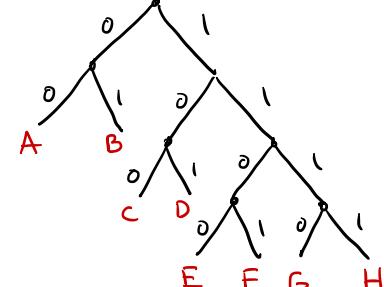
00|10|11101  
A D F

$$BPC = \sum_i l_i \cdot P_i$$

$$= 2 \cdot \frac{1}{4} + 2 \cdot \frac{1}{4} + 3 \cdot \frac{1}{8} + 3 \cdot \frac{1}{8} + 4 \cdot \frac{1}{16} + \dots + 4 \cdot \frac{1}{16}$$

$$= 2.75$$

Code Tree:



\* 8.33% savings over fixed code of length 3.