

Computer Networks: Theory, Modeling, and Analysis

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COM3510

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Lecture 1, 1

What is COM3510 about?

- Understand the basic principles of networking:
 - Description of existing networks, and networking mechanisms
 - Understanding of networks modeling and analysis tools
- Covers:
 - Terminology, layering concept, physical layer, data link layer, basics of queuing theory, medium access control (MAC), basics of routing, flow control (window/rate control)
- Does not cover: higher layers (applications, compression, TCP/IP programming, etc.)

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Course Outline

- Introduction to networking and to 7-layer OSI architecture
- Error detection, retransmissions strategies, and framing
- Introduction to Markov chains and queuing theory
- Multiple access schemes (MAC sublayer):
 - aloha, CSMA, ethernet, token rings, wireless, etc.
- Routing algorithms: Unicast and Multicast
 - Dijkstra, Bellman-Ford algorithms and RIP and OSPF protocols
- Packet Scheduling:
 - Max-Min Fairness, Weighted Fair Queuing.

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Lecture 1 Outline

- Networking Concepts and Terminology

- Protocols Layering

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Basic Terminology

- **Network:** set of nodes interconnected through communication links
- **Node:** host, router, switch
- **Link:** twisted-pair, coaxial cable, optical fiber, wireless
- **Protocol:** set of rules and conventions used between peer entities to communicate
- **Message:** sequence of bits/application level (e.g., email, document)
- **Packet:** messages are broken into packets that can be transmitted between network nodes
- **Session:** transaction consisting of a sequence of message exchanges

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Types of Networks

- Several taxonomies exist...
- Transmission technology based differentiation:
 - Point-to-point networks
 - Broadcast networks (generally small area: ethernet)
- Scale based differentiation:
 - Local Area Network (LAN): privately owned networks, up to few miles in size (e.g., ethernet)
 - Metropolitan Area Networks (MAN): larger than LANs, may cover a city (e.g., IEEE802.6 DQDB)
 - Wide Area Network (WAN): covers a large geographical area (e.g., country, continent)

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Performance Issues

- Packets arrival modeling assumed (simplicity) either:
 - *Poisson Process* (data packets)
 - *ON/OFF flow* model (for digitized voice)
- Common performance metrics:
 - latency and throughput,
 - delay-jitter, and rate-jitter (maximum - minimum)

Latency = *PropagationDelay* + *TransmissionDelay* + *QueuingDelay*

PropagationDelay = *Distance/SpeedOfLight* (independent of message size)

TransmissionDelay = *MessageSize/Bandwidth*

QueuingDelay = delay due to time spent waiting in queues

Sessions Characteristics

- Message arrival rate:
 - Poisson arrivals, deterministic arrivals, uniformly distributed
- Session holding time: duration of the session
- Expected message length and length distribution
- Allowable delay: depends on the application type
 - Interactive, real-time, email, voice, video-on-demand
- Reliability: depends on the application (FER)
 - Voice/video (tolerant), email (strict), sensing etc.
- Message & packet ordering: depends on the application
 - Databases, email, etc.

Types of Applications

- Interactive terminal and computer sessions:
 - Small packet length, small delay, high reliability
- File transfer:
 - High packet length, high delay, high reliability
- High resolution graphics:
 - High packet length, small delay, low arrival rate
- Voice application:
 - Small packet length, small delay, small reliability, high arrival rate

Sessions Transmission Paradigms

- Circuit Switching
- Store-and-Forward Switching

Circuit Switching

- On session establishment a path from source to destination is selected. Resources are allocated over all the links of the path. Route does not change during session life.
- Links can be shared by different sessions through mechanisms such as *time-division multiplexing* (TDM) or *frequency-division multiplexing* (FDM).
- For any link: the sum of rates of sessions using the link is at most the bandwidth of the link.
- Example: telephone networks.
- Guarantees: rate and in order packets delivery.

Store-and-Forward Switching

- Links are shared on a “*demand basis*” vs. fixed allocation
- Packets wait in a queue before being transmitted

Inefficiency of Circuit Switching

- Broadly speaking: when a session is idle, the reserved resources are lost.
- Simple analysis of session s :
 - λ : message arrival rate; P: propagation delay
 - $1/\lambda$: inter-arrival time between messages Q: queuing delay
 - \bar{X} : expected transmission time
 - \bar{L} : expected message length \bar{L}
 - Bit rate allocated to s : $r_s = \frac{\bar{L}}{\bar{X}}$
 - Link utilization = $\lambda\bar{X}$ $\bar{X} + P + Q \leq T \Rightarrow \lambda\bar{X} < \lambda T$
 - $T < 0.01$ for interactive app. *if $\lambda T \ll 1 \Rightarrow$ utilization – is – poor*

Circuit Switching vs. Store-and-Forward Switching

- Advantage of Store-and-Forward over Circuit Switching:
 - Network utilization is better (each link is utilized when there is some traffic)
 - Lower delays
- Drawbacks of Store-and-Forward:
 - Necessity of having control flow mechanisms to avoid buffer congestion and maintain acceptable delays. This is generally achieved through some feed-back to the senders.

More Taxonomy

- Message store-and-forward
- Packet store-and-forward
- Packet switching = store-and-forward switching

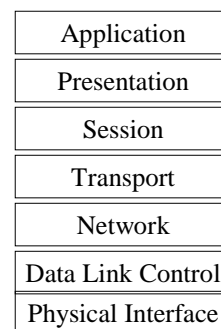
- Virtual circuit switching = packet switching + fixed path
- Dynamic routing => each packet finds its own path

Layering

- Layering: is a form of “*hierarchical modularity*”
- The role of each layer is to provide services for higher layers (abstracting lower layers)
- What is important is the functional relation between Inputs and Outputs
- Each layer has an interface to higher/lower layers (constituted of service access points: SAP)
- Entities at the same layer on different nodes are called *peer-entities*. They communicate through protocols
- A set of consecutive layers is called a *Protocol Stack*
- Drawback of layering: may hide important information (e.g., TCP over wireless links)

OSI Reference Model

- Open systems Interconnect (OSI) defined a reference for a layered architecture of data networks
- Existing protocol stacks (e.g., TCP/IP) are quite different from the OSI RM but it is still an interesting conceptual model because of its clean structure



Physical Layer

- Function: provides a “*virtual bit pipe*”
- How: maps bits into electrical/electromagnetic signals appropriate for the channel
- The physical layer module is called a *modem* (modulator/demodulator)
- Important issues:
 - Timing: synchronous, intermittent synchronous, asynchronous (characters)
 - Interfacing the physical layer and DLC (e.g., RS-232, X.21)

Data Link Control Layer (DLC)

- Receives packets from the network layer and transforms them into bits transmitted by the physical layer. *Generally* guarantees order and correctness.
- Mechanisms of the DLC:
 - Framing: header, trailer to separate packets, detect errors...
 - Multiple access schemes: when the link is shared by several nodes there is a need for addressing and controlling the access (this entity is called MAC sublayer)
 - Error detection and retransmission (LLC sublayer)

Network Layer

- Provides naming/addressing, routing, flow control, and scheduling/queuing in a multi-hop network
- Makes decisions based on packet header (e.g., destination address) and module stored information (e.g., routing tables)
- General comment: each layer looks only at its corresponding header (here packet header)
- Routing is different on virtual circuit networks than on datagram networks

Datagram vs. Virtual Circuits

- In virtual circuits after an initialization phase all packets follow the same path. We generally assume that packets are delivered once and only once, and in order.
- In a datagram network packets are routed individually. They may be lost or delivered out-of sequence.
- Sometimes referred to as: *connection-oriented service* and *connectionless service*.

Flow control & Congestion

- Flow control avoids sending data faster than the destination can absorb
- Congestion control avoids sending data faster than the network can handle.
- Requires:
 - efficient feedback mechanism, buffer management, route load-balancing
- In a connectionless service it is not easy to negotiate an equitable service between users

Transport Layer

- Provides a reliable mean to transmit messages between two end-nodes through:
 - Messages fragmentation into packets
 - Packets reassembly in original order
 - Sessions multiplexing and splitting
 - Retransmission of lost packets
 - end-to-end flow control and congestion control

Session Layer

- Was intended to handle the interaction between two end points in setting up a session:
 - Two connections
 - Service location (e.g., would achieve load sharing)
 - Control of access rights
- In many networks these functionalities are inexistent or spread over other layers

Presentation Layer

- Was intended to provide data encryption, data compression, and code/types conversion from one architecture to another.

Application Layer

- What's left over...
- Examples: WWW, Email, Telnet, ...

Simple Distributed Problem

- Three armies: two on the same side (A), one enemy B.
- If the two A armies attack they win, but if only one attacks it loses.
- Communication between A armies goes through an area controlled by B army.
- Can A armies find a strategy to defeat B.

Summary

- Networking terminology and basic concepts
- Layering in networking
- 7-layers OSI reference model
- Brief introduction of each layer