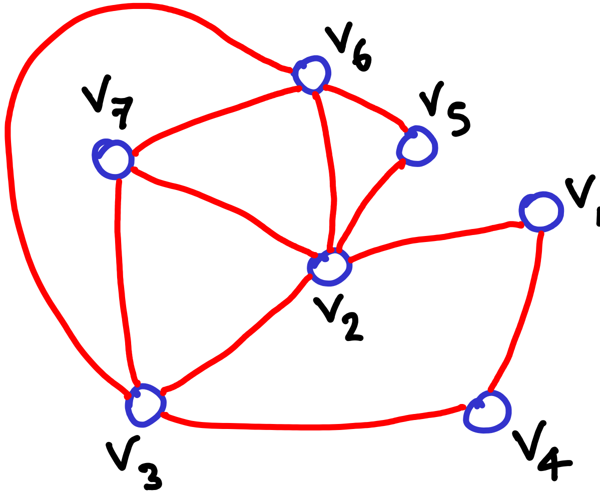
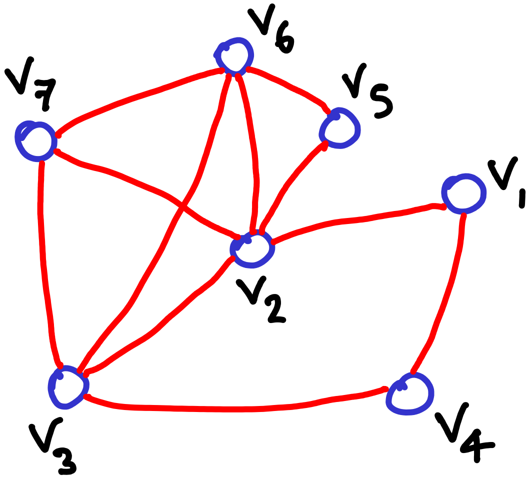
PLANE GRAPH

no crossings

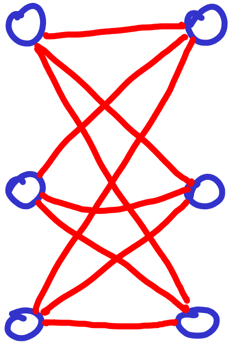




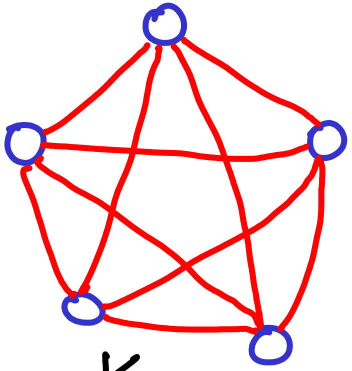
PLANAR GRAPH

can redraw without crossings

Non-planar graphs
(can't redraw)

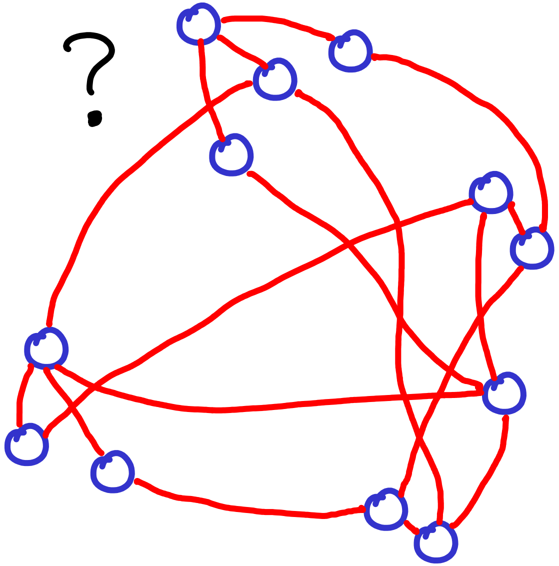


$K_{3,3}$



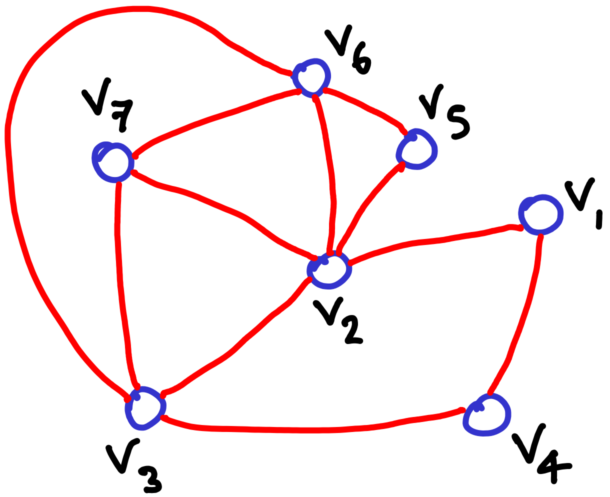
K_5

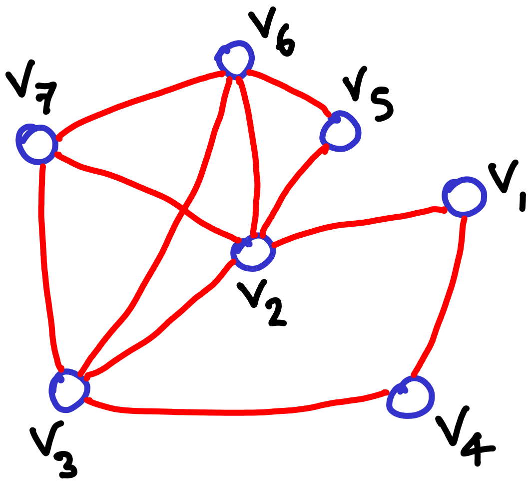
?



PLANE GRAPH

no crossings



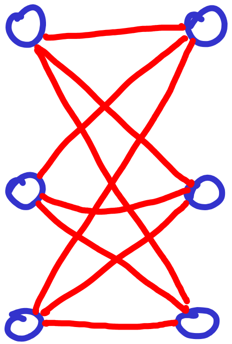


PLANAR GRAPH

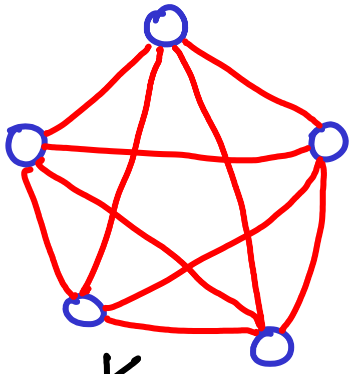
can redraw without crossings

Non-planar graphs

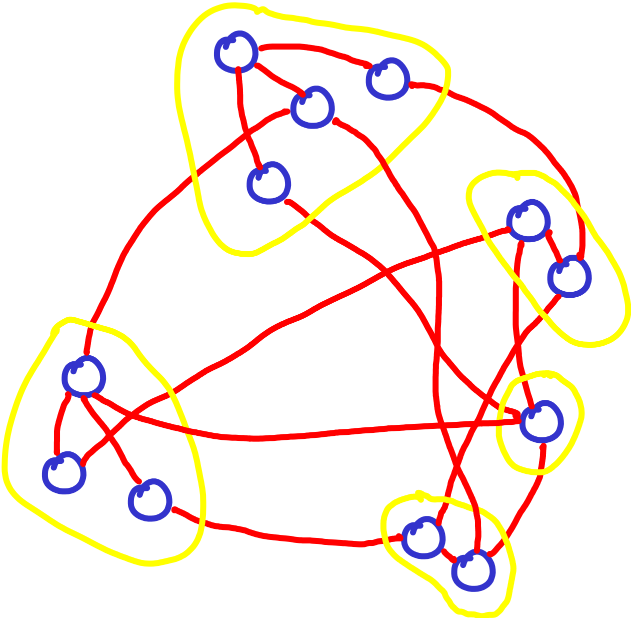
(can't redraw)



$K_{3,3}$

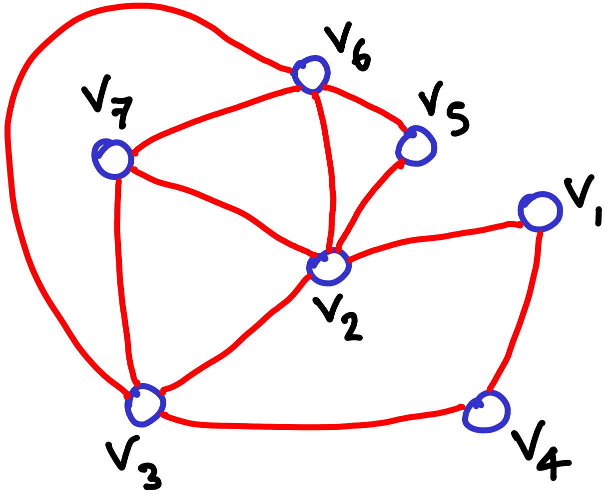


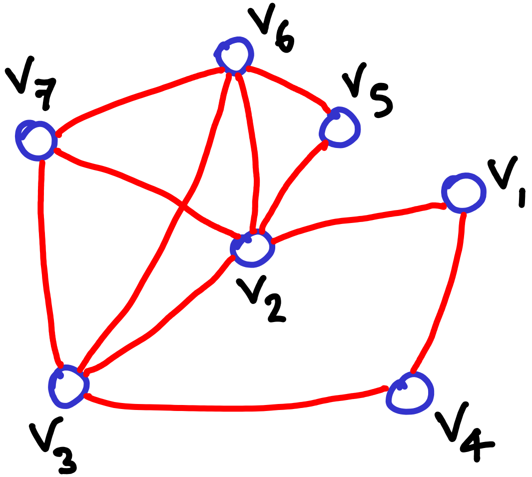
K_5



PLANE GRAPH

no crossings

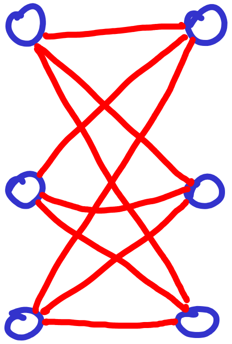




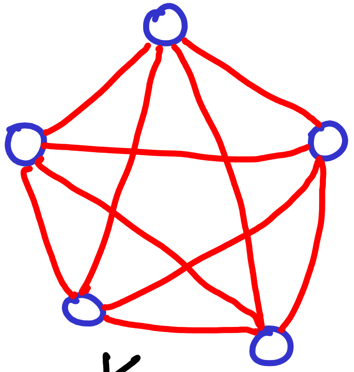
PLANAR GRAPH

can redraw without crossings

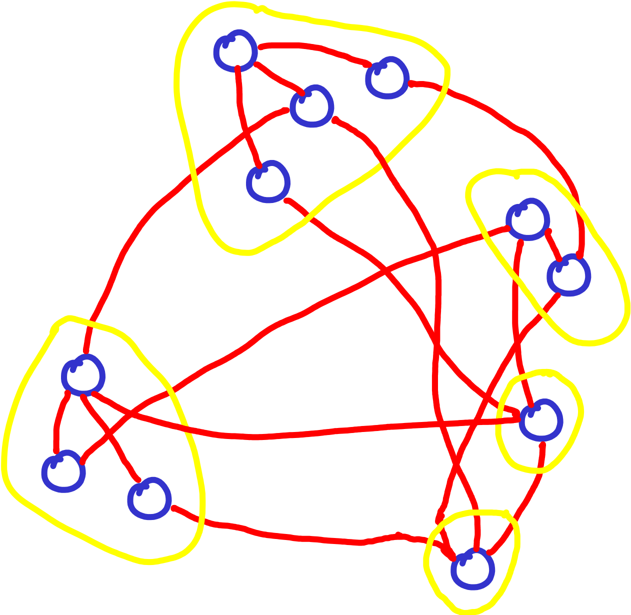
Non-planar graphs (can't redraw)



$K_{3,3}$

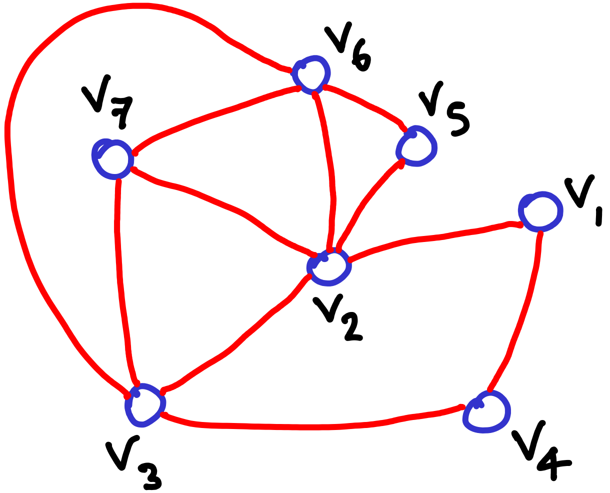


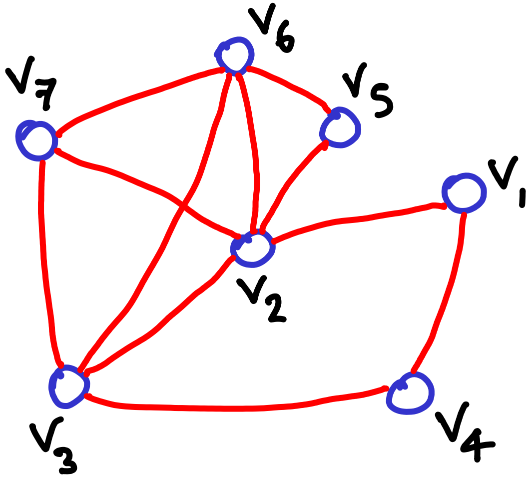
K_5



PLANE GRAPH

no crossings



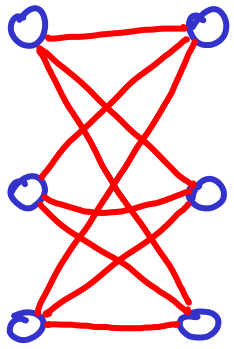


PLANAR GRAPH

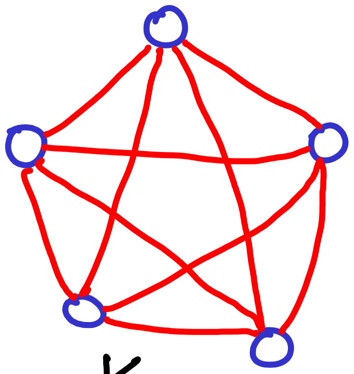
can redraw without crossings

Non-planar graphs

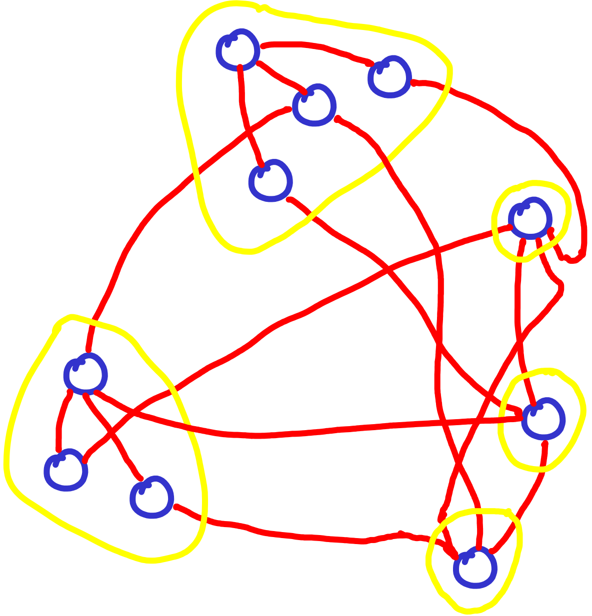
(can't redraw)



$K_{3,3}$

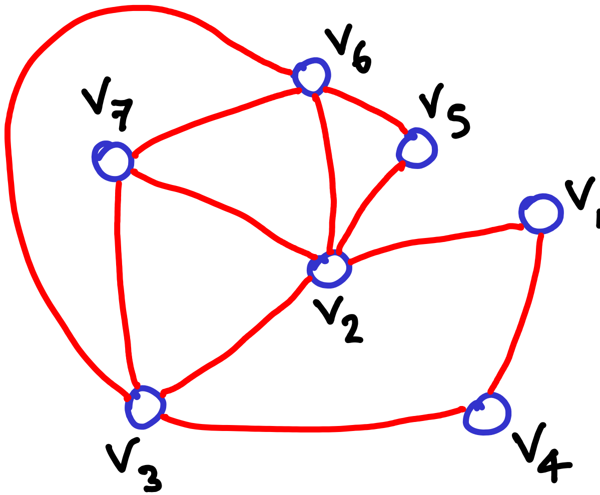


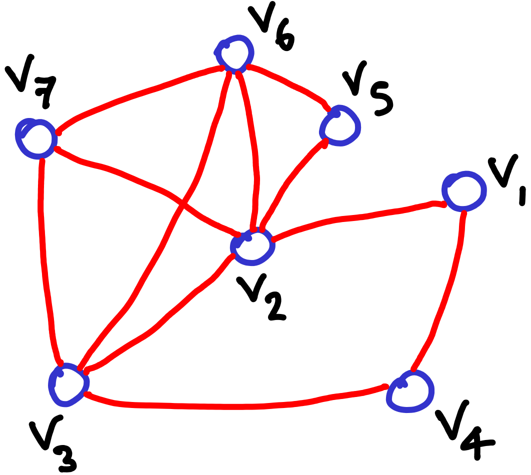
K_5



PLANE GRAPH

no crossings



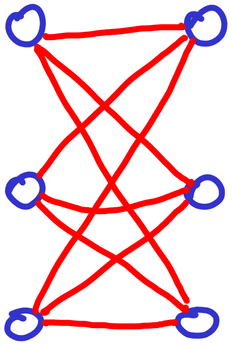


PLANAR GRAPH

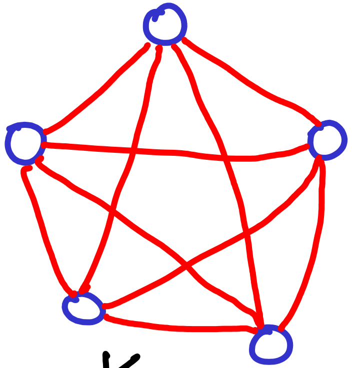
can redraw without crossings

obtained by successive contractions

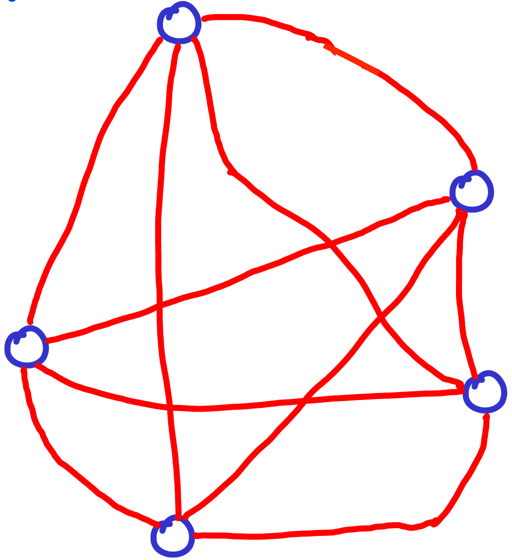
Non-planar graphs (can't redraw)



$K_{3,3}$



K_5



A graph is non-planar if and only if it "contains" a K3,3 or K5

FYI

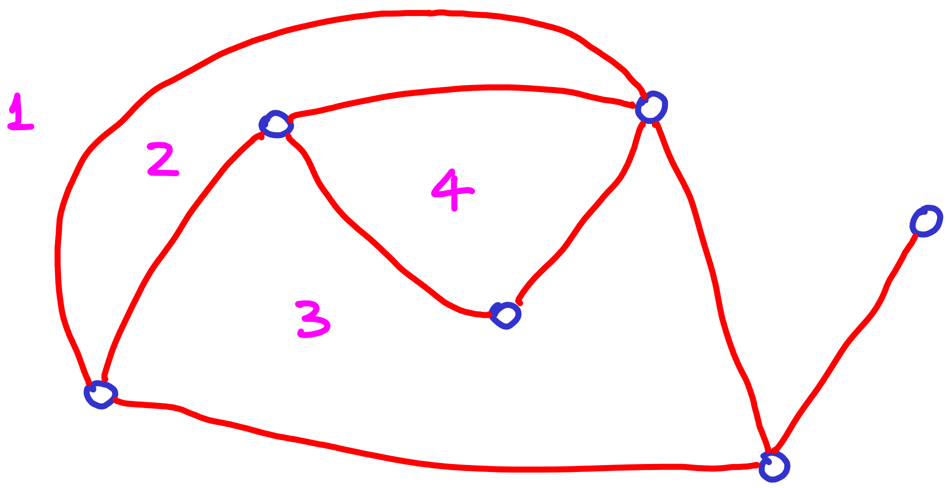
Every planar graph can be drawn without crossings.
In fact the edges can be drawn straight as well.

$G = (V, E)$

$V = 6$

$E = 8$

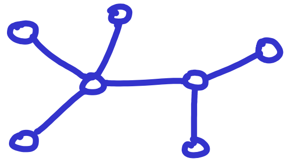
$F = \underbrace{\#\text{faces}}_{\text{disjoint regions}} = 4$



disjoint regions, one of which is unbounded

EULER FORMULA for planar connected graphs: $\underline{V - E + F = 2}$

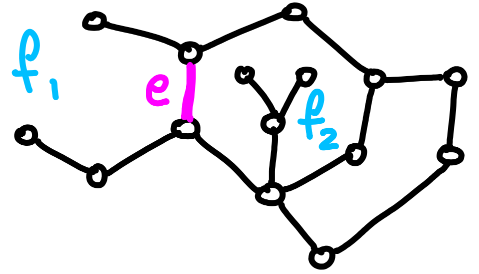
Proof by induction on number of faces:



Base case $\to$ $F = 1$ $\to$ G is a tree $\to$ $V = E + 1$

so $(E+1) - E + 1 = 2$ ✓

Given $G = (V, E)$ w/ $F > 1$ faces,



remove an edge e between 2 faces, f_1 & f_2.

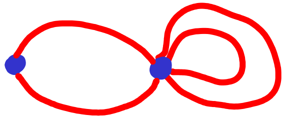
Either f_1 or f_2 is a bounded face,

so e is on a cycle (e is not a cut edge)

$\hookrightarrow$ $G - e$ is connected & f_1, f_2 merge:

$$\overbrace{V - (E - 1) + (F - 1) = 2}^{\text{hypothesis}}$$

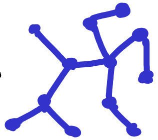
$\hookrightarrow$ $V - E + F = 2$ ✓

Note that this also holds for multigraphs 

Euler formula $V - E + F = 2$

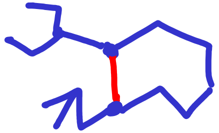
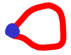
$V - E + F = 2$ applies to any connected planar graph (in fact, to convex polyhedra by projection)

Induction on faces:

$F = 1$:  : tree. $E = V - 1$

$F > 1$:

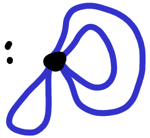
Remove an edge between 2 faces.

Remains connected.

$F \to F-1 \quad E \to E-1$

Induction on vertices

$V = 1$:  : only loops $F = E + 1$

$V > 1$:

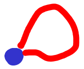
Contract edge $x \neq y$ 

$V \to V-1 \quad E \to E-1$

Induction on edges

$E = 0$:  : one vertex one face

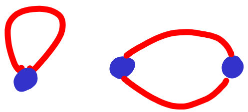
$E \geq 1$: if $x \neq y$  contract as before

else  remove as before

$E \to E-1$ & $F \to F-1$ OR $V \to V-1$

use the Euler formula $V - E + F = 2$

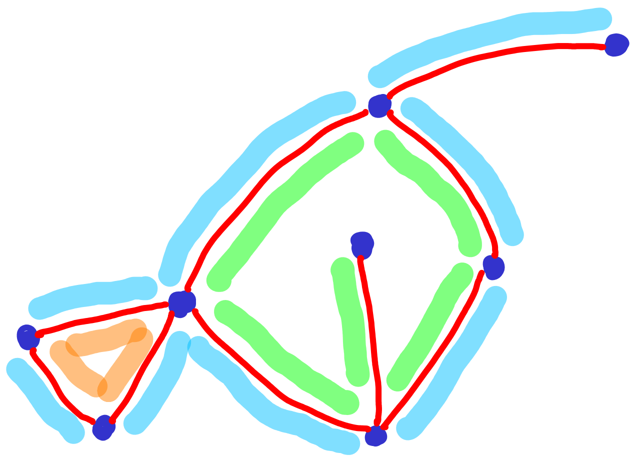
to show that a connected plane graph has $E \leq 3V - 6$ for $V > 3$

Not allowed: 

use the Euler formula $V - E + F = 2$

to show that a connected plane graph has $E \leq 3V - 6$

Every edge belongs to 1 or 2 faces $\sum_{\text{all faces}} e \leq 2E$

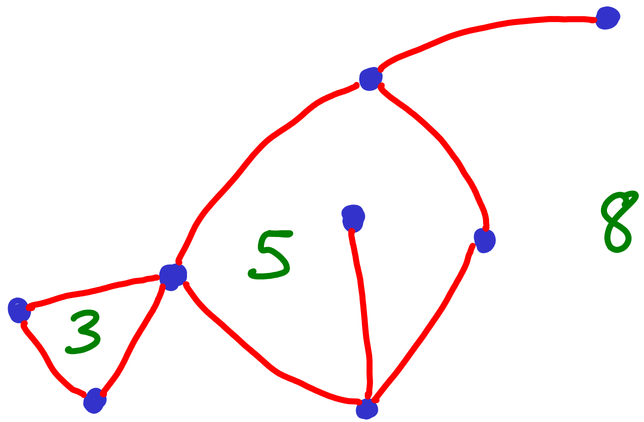


use the Euler formula $V-E+F=2$

to show that a connected plane graph has $E \leq 3V-6$

Every edge belongs to 1 or 2 faces $\quad \sum_{\text{all faces}} e \leq 2E$

Every face has ≥ 3 edges (for $V>3$) $\quad \sum_{\text{all faces}} e \geq 3F$



use the Euler formula $V - E + F = 2$

to show that a connected plane graph has $E \leq 3V - 6$

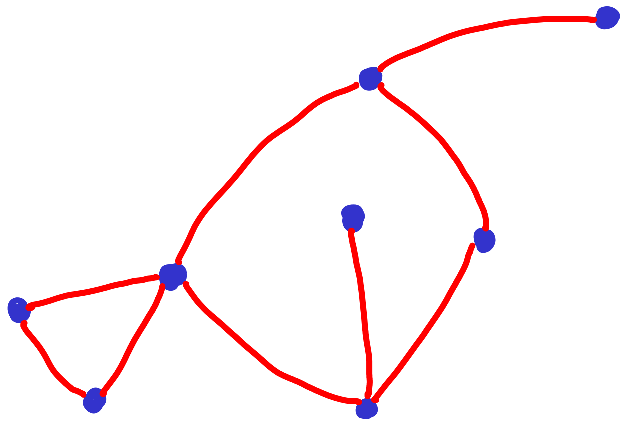
Every edge belongs to 1 or 2 faces $\quad \sum_{\text{all faces}} e \leq 2E$

Every face has ≥ 3 edges (for $V > 3$) $\quad \sum_{\text{all faces}} e \geq 3F$

$2E \geq 3F$

$$E - F = V - 2$$
$$E - \frac{2E}{3} \leq V - 2$$
$$E \leq 3V - 6$$

$$E - F = V - 2$$
$$\frac{3F}{2} - F \leq V - 2$$
$$F \leq 2V - 4$$

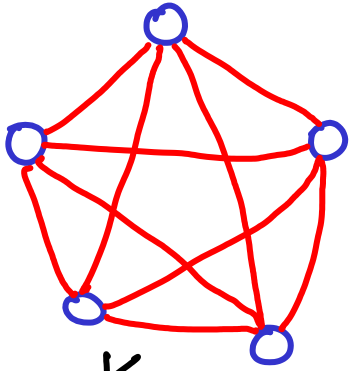


$E \leq 3V-6$

$10 \leq 15-6$

!!!

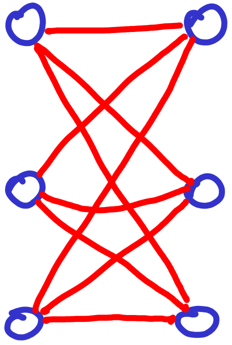
not planar



K_5

$E \leq 3V-6$

$9 \leq 18-6$ ok!



$K_{3,3}$

not iff

All planar graphs have $E \leq 3V-6$

Some non-planar graphs can too

$V - E + F = 2$

What if G has no triangles?

Every edge belongs to 1 or 2 faces $\qquad \sum_{\text{all faces}} e \leq 2E$

Every face has $\overset{\geq 4}{\cancel{\geq 3}}$ edges (for $V > 4$) $\qquad \sum_{\text{all faces}} e \geq 4F$

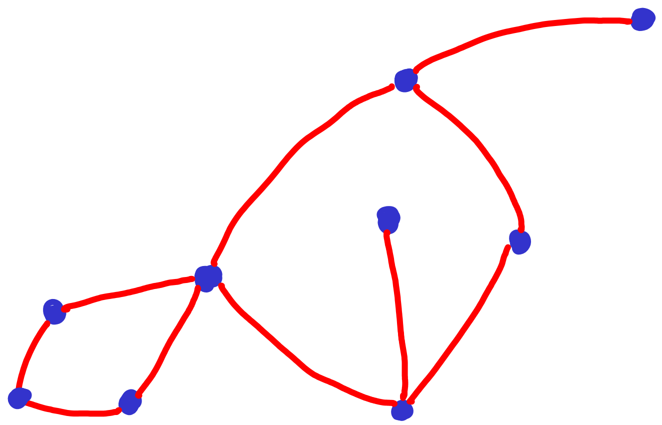
$E \geq 2F$

$E - F = V - 2$

$E - \frac{E}{2} \leq V - 2$

$E \leq 2V - 4$

instead of $\leq 3V - 6$



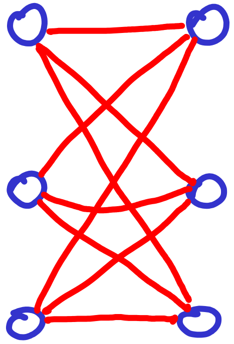
for triangle free: $E \leq 2V - 4$

$9 \leq 2 \cdot 6 - 4$

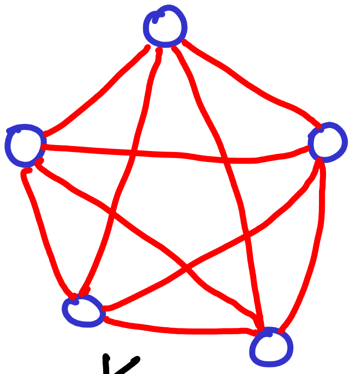
!!!

NOT PLANAR

$V=6,\ E=9$

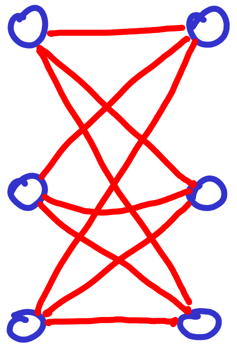


$K_{3,3}$



K_5

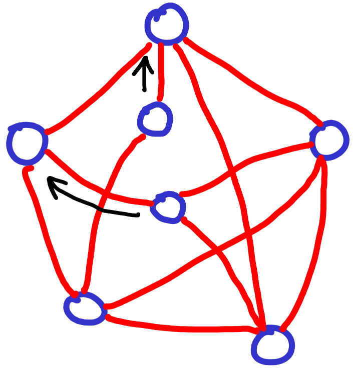
⟵ non-planar ⟶



$K_{3,3}$

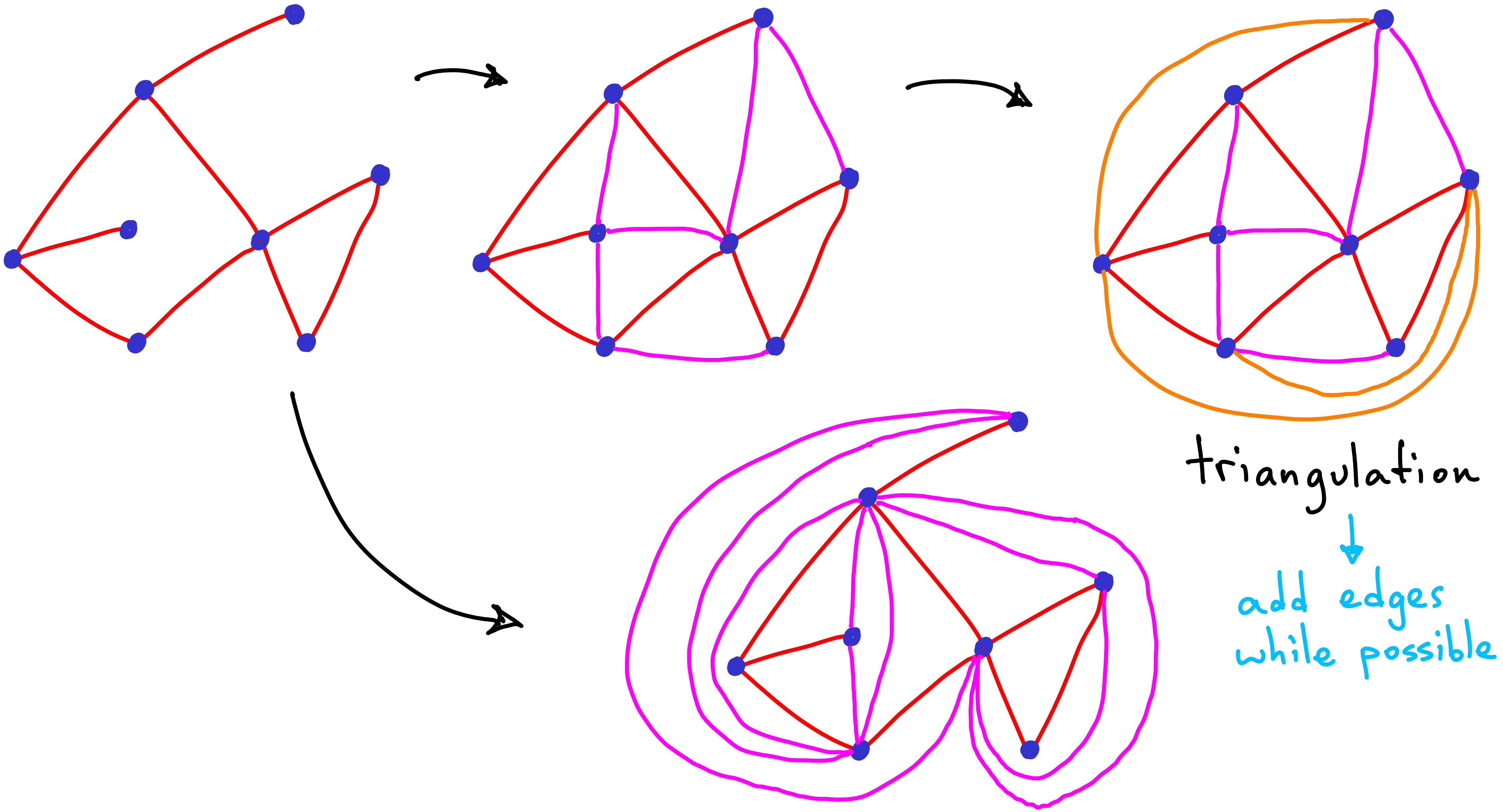
It turns out that every non-planar graph "contains" one of these two shapes.

e.g:



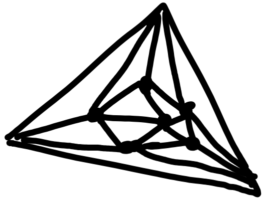
see links

(example was shown, with contractions)


triangulation
add edges while possible

$\underline{\underline{E = 3V - 6}}$

Why?



{ Assume outer face is a triangle

$\boxed{V - E + F = 2}$

Every edge belongs to ~~1 or~~ 2 faces $\qquad \sum_{\text{all faces}} e \overset{=}{\cancel{\leq}} 2E$

Every face has $\overset{=}{\cancel{\geq}}$ 3 edges (for $V > 3$) $\qquad \sum_{\text{all faces}} e \overset{=}{\cancel{\geq}} 3F$

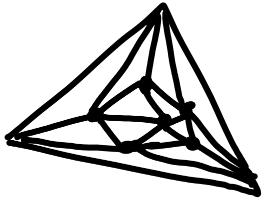
$2E \overset{=}{\cancel{\geq}} 3F$

$$E - F = V - 2$$

$$E - \frac{2E}{3} = V - 2$$

$$E = 3V - 6$$

$E = 3V - 6$



What is the average degree of a triangulation?

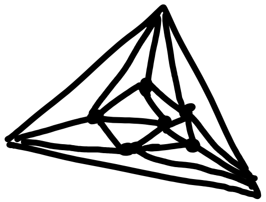
$$\frac{1}{V} \cdot \sum_{i=1}^{V} d(v_i) = \frac{1}{V} \cdot 2E = \frac{6V-12}{V} < 6$$

↪ Every triangulation has a vertex w/ degree ≤ 5

↪ Immediately applies to any planar graph

(fewer edges)

$E = 3V-6$



What is the average degree of a triangulation?

$$\frac{1}{V}\cdot\sum_{i=1}^{V} d(v_i) = \frac{1}{V}\cdot 2E = \frac{6V-12}{V} < 6$$

↳ Every planar graph has a vertex w/ degree ≤ 5

Can we find many low-degree vertices? → not if "low" = 5.
what if "low" = 8 ?

Say you had $\geq \frac{V}{2}$ vertices w/ degree ≥ 9

↳ $\sum_{d \geq 9} d(v_i) \geq 9 \cdot \frac{V}{2}$

sum degrees of $V/2$ of them

all other vertices have degree ≥ 3

↳ $\sum d(v_i) \geq 3 \cdot \frac{V}{2}$

$\Sigma + \Sigma \geq 6V$

contradiction

$\Downarrow$

always have $\geq \frac{V}{2}$ w/ deg ≤ 8